Maximum Weight Independent Sets in Odd-Hole-Free Graphs Without Dart or Without Bull

Andreas Brandstädt[*]   Raffaele Mosca[†]

Abstract

The Maximum Weight Independent Set (MWIS) Problem on graphs with vertex weights asks for a set of pairwise nonadjacent vertices of maximum total weight. Being one of the most investigated and most important problems on graphs, it is well known to be NP-complete and hard to approximate. The complexity of MWIS is open for hole-free graphs (i.e., graphs without induced subgraphs isomorphic to a chordless cycle of length at least five). By applying clique separator decomposition as well as modular decomposition, we obtain polynomial time solutions of MWIS for odd-hole- and dart-free graphs as well as for odd-hole- and bull-free graphs (*dart* and *bull* have five vertices, say a, b, c, d, e, and dart has edges ab, ac, ad, bd, cd, de, while bull has edges ab, bc, cd, be, ce). If the graphs are hole-free instead of odd-hole-free then stronger structural results and better time bounds are obtained.

Keywords: Maximum weight independent set; clique separators; modular decomposition; polynomial time algorithm; hole-free graphs; dart-free graphs; bull-free graphs.

1 Introduction

The Maximum Weight Independent Set (MWIS) Problem on a given finite undirected simple graph with vertex weights asks for a set of pairwise nonadjacent vertices of maximum total weight. Being one of the most investigated and most important problems on graphs, it is well known to be NP-complete and hard to approximate. It is solvable in polynomial time on various graph classes while it remains NP-complete on some others. Its complexity is open for hole-free graphs (i.e., graphs without induced subgraphs isomorphic to a chordless cycle of length at least five). Recently, the following subclasses of hole-free graphs were studied:

(i) hole-free graphs without induced diamond and in general without induced paraglider [6];

(ii) hole-free graphs without induced co-chair [5];

(iii) hole-free graphs without induced dart [1].

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[*]Institut für Informatik, Universität Rostock, D-18051 Rostock, Germany. E-mail: ab@informatik.uni-rostock.de

[†]Dipartimento di Economia, Universitá degli Studi "G. d'Annunzio", Pescara 65121, Italy. E-mail: r.mosca@unich.it

In [6], it was shown that the atoms of the graphs in the class (i) are weakly chordal or specific graphs, and in [5], it was shown that the prime atoms of the graphs in the class (ii) are nearly weakly chordal, from which polynomial time algorithms for MWIS on these graphs follow.

Using the approach of [5, 6], Basavaraju, Chandran and Karthick in [1] showed that the MWIS problem can be solved in polynomial time for hole- and dart-free graphs by reducing it first to hole-, dart-, and gem-free graphs (which are hole- and paraglider-free). We extend previous results and show:

(i) odd-hole- and dart-free graphs are nearly perfect, and (hole, dart)-free atoms are nearly weakly chordal.

(ii) odd-hole- and bull-free prime graphs are nearly perfect, and (hole, bull)-free prime graphs are nearly weakly chordal.

(iii) hole- and dart-free graphs as well as hole- and bull-free graphs have nice structure properties; MWIS for hole- and bull-free graphs can be solved in time $\mathcal{O}(n^5)$.

(iv) MWIS for P_5- and bull-free graphs can be solved in time $\mathcal{O}(nm)$.

The results in (i) and (ii) are based on the Strong Perfect Graph Theorem and imply that the MWIS problem is solvable in polynomial time for odd-hole- and dart-free graphs as well as for odd-hole- and bull-free graphs.

Actually, Lemma 1 in [5] implies that also (odd-hole,co-chair)-free prime atoms are nearly perfect (which implies polynomial time for the MWIS problem on (odd-hole,co-chair)-free graphs). This fact was not explicitly mentioned in [5].

2   Some Basic Notions and Results

2.1   Basic Notions

For any missing notation or reference let us refer to [9]. Let $G = (V, E)$ be a finite undirected graph which is simple (i.e., without self-loops and multiple edges) with vertex set V and edge set E. Let $|V| = n$ and $|E| = m$. For a vertex $v \in V$, let $N(v) := \{u \mid uv \in E\}$ denote the *open neighborhood* of v and let $N[v] := N(v) \cup \{v\}$ denote the *closed neighborhood* of v. We also say that for vertices $u, v \in V$, u and v *see each other* (*miss each other*, respectively) if $uv \in E$ ($uv \notin E$, respectively). Let $A(v) := V \setminus N[v]$ denote the *anti-neighborhood* of v.

For any vertex set $U \subseteq V$, let $N(U) := \bigcup_{u \in U} N(u) \setminus U$ and $N[U] := N(U) \cup U$ as well as $A(U) := V \setminus N[U]$. For any nonempty vertex subset $U \subseteq V$ with $A(U) \neq \emptyset$, let $U^+ := N(U) \cap N(A(U))$ (the set of *contact vertices* of U and $A(U)$).

For any vertex set $U \subseteq V$ let $G[U]$ be the subgraph of G induced by U. Let $\overline{G} = (V, \overline{E})$ denote the *complement graph of G*, also denoted as co-G. Pairs $xy \in \overline{E}$ are also called *co-edges* of G.

For any disjoint vertex sets U, W of V, let us say that U has a *join* (a *co-join*, respectively) to W if each vertex of U sees each vertex of W (misses each vertex of W, respectively). If $U = \{u\}$ and U has a join to W, then let us say that u *dominates* W or is *universal for* W.
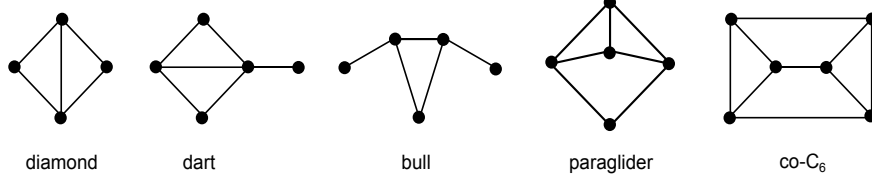
Figure 1: diamond, dart, bull, paraglider, and co-C_6

P_k is the induced path with k vertices and $k-1$ edges. C_k is the induced cycle with k vertices and k edges. A *hole* is C_k with $k \geq 5$. An *odd hole* is C_{2k+1} with $k \geq 2$. An *anti-hole* is the complement graph of a hole. A *diamond* (or $K_4 - e$) is formed by vertices a, b, c, d, and edges ab, ac, ad, bd, cd. A *gem* is a one-vertex extension of a diamond, which can be obtained by adding a dominating vertex to a P_4. A *dart* has five vertices a, b, c, d, e, and edges ab, ac, ad, bd, cd, de, i.e., it consists of a diamond plus a degree-one vertex being adjacent to one of the degree-3 vertices of the diamond. A *bull* has five vertices a, b, c, d, e and edges ab, bc, cd, be, ce. See Figure 1 for most of these specific graphs.

For any graph H, let us say that G is *H-free* if G contains no induced subgraph isomorphic to H. A class of graphs is *hereditary* if it is closed under taking induced subgraphs. A graph is *chordal* it is hole-free and C_4-free. A graph is *weakly chordal* if it is hole- and anti-hole-free. *Perfect graphs* play a crucial role in algorithmic graph theory, and the Strong Perfect Graph Theorem (see Theorem 3) characterized them in terms of forbidden odd holes and odd anti-holes. It is well known that chordal graphs are weakly chordal, and weakly chordal graphs are perfect (see e.g. [9] for this and the many facets of these graphs).

An *independent set* of G is a subset of pairwise non-adjacent vertices of G. A *clique* of G is a set of pairwise adjacent vertices of G.

2.2 Techniques for the Maximum Weight Independent Set Problem

The Maximum Independent Set (MIS) Problem asks for an independent vertex set of maximum cardinality. If w is a real-valued function on V then the Maximum Weight Independent Set (MWIS) Problem asks for an independent set of maximum total weight. Let $\alpha_w(G)$ denote the maximum weight of an independent vertex set in G.

A subset $U \subseteq V$ is a *cutset* (or *separator*) in G if $G[V \setminus U]$ has more connected components than G. A *clique cutset* is a cutset which is a clique. An *atom* in G is an induced subgraph of G without clique cutset. More generally, a graph is an *atom* if it has no clique cutset. A famous divide-and-conquer approach by using clique separators is described in [28, 29]. A consequence is the following:

Theorem 1 ([28, 29]) *If for a hereditary graph class $\mathcal{C}$, the MWIS problem is solvable in polynomial time for the atoms of $\mathcal{C}$ then the MWIS problem is solvable in polynomial time on graph class $\mathcal{C}$.*

Let Π denote a (hereditary) graph property. A graph $G = (V, E)$ is *nearly* Π if for all $v \in V$, the subgraph $G[A(v)]$ has property Π. For short, we say that for every vertex,

3

the anti-neighborhood has property Π. For example, G is *nearly weakly chordal* if the anti-neighborhood of every vertex is weakly chordal. Obviously the following holds:

Observation 1 $\alpha_w(G) = \max\{w(v) + \alpha_w(G[A(v)]) \mid v \in V\}$

Thus, we obtain:

Corollary 1 *Whenever MWIS can be solved in time T on a hereditary graph class with property Π, it can be solved in time $|V| \cdot T$ on nearly Π graphs.*

For example, since the MWIS problem can be solved in time $\mathcal{O}(n^4)$ for weakly chordal graphs [27], it can be solved in time $\mathcal{O}(n^5)$ for nearly weakly chordal graphs.

In Section 4 we make use of modular decomposition. We say that a vertex z *distinguishes two vertices* x, y if z sees x and misses y. A subset U of vertices is a *module in G* if no vertex $z \in V \setminus U$ distinguishes two vertices $x, y \in U$. A module is *trivial* if it is either the empty set or V or one-elementary. A graph is *prime* if all its modules are trivial. We will use primality for solving MWIS on ((odd-)hole,bull)-free graphs. It is well known that MWIS can be solved in time T bottom-up along the modular decomposition tree for a hereditary graph class $\mathcal{C}$ if MWIS can be solved in time T for the prime graphs of $\mathcal{C}$; the modular decomposition tree of a given graph can be determined in linear time [24].

Theorem 2 *If for a hereditary graph class $\mathcal{C}$, the MWIS problem is solvable in time T, $T \geq m$, for the prime graphs of $\mathcal{C}$ then the MWIS problem is solvable in time $\mathcal{O}(T)$ on graph class $\mathcal{C}$.*

2.3 Basic Results on Some Classes of Perfect and Other Graphs

Subsequently, we use the following results and facts:

Theorem 3 (Strong Perfect Graph Theorem [12]) *A graph is perfect if and only if it is odd-hole-free and odd-anti-hole-free.*

Theorem 4 *The following graph classes can be recognized in polynomial time:*

(i) *weakly chordal graphs [2, 23, 27]*

(ii) *hole-free graphs [19, 26]*

(iii) *perfect graphs [15].*

Recognition of odd-hole-free graphs is open, though recognition of odd-hole-free graphs with cliques of bounded size can be done in polynomial time [14].

Theorem 5 *The MWIS problem for can be solved in polynomial time for the following graph classes:*

(i) *weakly chordal graphs [27]; the time bound is $\mathcal{O}(n^4)$.*

(ii) *perfect graphs [22].*

(iii) *hole-free graphs with no banner [8], with no paraglider [6], with no co-chair [5] and with no dart [1].*

The MWIS problem for hole-free graphs (and for odd-hole-free graphs) is open; it is open even for the subclass of (P_5, C_5)-free graphs.

Finally let us mention a recent paper [10] introducing many structural properties for bull-free graphs (see also [18] for perfect bull-free graphs) and the papers [7, 17] focussing on efficiently solving MWIS for (bull,chair)-free graphs.

3 Structure and MWIS for (Odd-)Hole- and Dart-Free Graphs

In this section, we show that odd-hole- and dart-free graphs are nearly perfect and that hole- and dart-free atoms are nearly weakly chordal. This is the main result of the section and implies polynomial time for MWIS on both graph classes (with better time bound for hole- and dart-free graphs). We first collect some properties of dart-free graphs and deal with (odd-hole,dart)-free graphs..

3.1 Structure and MWIS for (Odd-Hole,Dart)-Free Graphs

Proposition 1 *Let $G = (V, E)$ be a dart-free graph and let $U \subset V$. If $u \in U^+$ and vertices a, b, c induce a P_3 in U then u does not see all three vertices a, b, c.*

Proof. Otherwise, since $u \in U^+$ sees a vertex $v \in A(U)$ (recall that $A(U) \neq \emptyset$), such vertices a, b, c together with u and v would induce a dart. □

Subsequently, when dealing with cycles of length k as well as their complements, index arithmetic is done modulo k.

Lemma 1 *Let $G = (V, E)$ be a dart-free graph containing an anti-hole $\overline{C_k}$, $k \geq 7$, say H, with vertices $v_1, \ldots, v_k$ and co-edges $v_i v_{i+1}$ for $i = 1, \ldots, k$, such that $H^+ \neq \emptyset$, and let $x \in H^+$. Then the following hold:*

 (i) If x sees vertex v_i for some index i, then x sees v_{i-2} and v_{i+2}.

 (ii) k is even, and either $N_H(x) = \{v_1, v_3, \ldots, v_{k-1}\}$ or $N_H(x) = \{v_2, v_4, \ldots, v_k\}$.

Proof. (i): Assume without loss of generality that x sees v_1. First, we show that x sees either v_3 or v_{k-1}. Assume to the contrary that x sees neither v_3 nor v_{k-1}. Then to avoid that $x, v_1, v_3, v_{k-1}, v_4$ induce a dart, x sees v_4, and then by Proposition 1, x misses v_k (since $x \in H^+$). Similarly by symmetry one has that x sees v_{k-2}, and x misses v_2. If $k = 7$, then v_4 misses v_{k-2}, a contradiction to Proposition 1. If $k > 7$, then v_4 sees v_{k-2}, and then $v_{k-1}, v_4, x, v_k, v_{k-2}$ induce a dart, a contradiction. Then x sees either v_3 or v_{k-1}.

Then let us assume without loss of generality that x sees v_3, so by Proposition 1, x misses v_4 and v_k but then x sees v_{k-1} too, otherwise $v_k, v_3, x, v_1, v_{k-1}$ induce a dart.

(ii): By Proposition 1, no vertex of H^+ dominates H. Thus, if k is odd, then by an iterated application of statement (i), one has that each vertex $x \in H^+$ dominates H, a contradiction to Proposition 1; if k is even then by statement (i) the last part of statement (ii) follows. □

Corollary 2 *Dart-free graphs are nearly $\overline{C_{2k+1}}$-free for $k \geq 3$.*

Theorem 3, Lemma 1 and Corollary 2 imply the following result:

Theorem 6 (*Odd-hole,dart)-free graphs are nearly perfect.*

Proof. Let G be an (odd hole, dart)-free graph. To prove the assertion it is sufficient to recall Theorem 3 and to show that if G has a odd anti-hole H with at least 7 vertices, then $H^+ = \emptyset$ which follows by contradiction from Lemma 1 (ii). □

Theorem 5 (ii) implies:

Corollary 3 *The MWIS problem is solvable in polynomial time for (odd-hole,dart)-free graphs.*

3.2 Structure and MWIS for (Hole,Dart)-Free Graphs

For (hole,dart)-free graphs, some stronger properties can be shown if one additionally excludes clique cutsets:

Lemma 2 (*Hole, dart)-free atoms are nearly $\overline{C_k}$-free for $k \geq 7$.*

Proof. Assume to the contrary that there is a vertex v in G such that the anti-neighborhood $A(v)$ of v contains an induced $\overline{C_k}$, say H, for $k \geq 7$, with vertices $v_1, \ldots, v_k$ and co-edges $v_i v_{i+1}$ for $i = 1, \ldots, k$ (index arithmetic modulo k). Clearly $H^+ \neq \emptyset$ and H^+ is a cutset for G. Let Q_v be the connected component of the anti-neighborhood $A(H)$ of H containing v. Then since G is an atom, there exist two vertices, say $x, y \in H^+$, such that:

(a) x misses y, and

(b) both x and y see some vertex of Q_v.

Then by Lemma 1, let us distinguish between the following two cases, which are exhaustive by symmetry:

1. If $N_H(x) = \{v_1, v_3, \ldots, v_{k-1}\}$ and $N_H(y) = \{v_1, v_3, \ldots, v_{k-1}\}$, then x, y, v_1, v_3, v_4 induce a dart, a contradiction.

2. If $N_H(x) = \{v_1, v_3, \ldots, v_{k-1}\}$ and $N_H(y) = \{v_2, v_4, \ldots, v_k\}$, then x, v_5, v_2, y induce a P_4; on the other hand, let P be a shortest path form x to y in $G[Q_v \cup \{x, y\}]$; then the subgraph induced by P, x, y, v_1, v_2 is a hole, a contradiction.

This finally shows Lemma 2. □

The proof of the subsequent Lemma 3 is similar to the one of Theorem 2 in [5]; in particular, the parts which are exactly the same (i.e., those which did not require the assumption co-chair-freeness) are reported in their original form.

Lemma 3 (*Hole, dart)-free atoms are nearly $\overline{C_6}$-free.*

Proof. Assume to the contrary that there is a vertex v in G such that its anti-neighborhood $A(v)$ contains an induced $\overline{C_6}$, say A, with vertices $v_1, \ldots, v_6$ such that v_1, v_2, v_3 is a clique left(A), v_4, v_5, v_6 is a clique right(A), and $v_1 v_4$, $v_2 v_5$, and $v_3 v_6$ are the edges between left(A) and right(A) (the *matching edges of A*). Let A_i denote the neighbors of A which see exactly i vertices in A, $1 \leq i \leq 6$, and let A^+ denote the neighbors of A which see a vertex in the connected component Q_v of the anti-neighborhood of A containing v. Note that A^+ depends on v but for short (and in order to avoid confusion with anti-neighborhoods) we write A^+ instead of $A^+(v)$.

Let also $A_2^+(1,4)$ denote the vertices of A_2^+ which see exactly v_1 and v_4 in A and similarly for some other cases. We first collect some simple properties.

Claim 3.1 $A_4^+ = A_5^+ = A_6^+ = \emptyset$

Proof of Claim 3.1. Follows by Proposition 1. $\diamond$

Claim 3.2 *If $x \in A_2$ then x sees the two vertices of a matching edge in A.*

Proof of Claim 3.2. Assume not; then x is either adjacent to two vertices in left(A) (right(A) respectively), say x sees v_1, v_2, in which case v_4, v_3, v_2, v_1, x induce a dart, or x is adjacent to two nonadjacent vertices, say x sees v_1, v_5 in which case x, v_1, v_3, v_6, v_5 induce a C_5. $\diamond$

Claim 3.3 *If $x \in A_3^+$ then x has a join to $left(A)$ or x has a join to $right(A)$.*

Proof of Claim 3.3. Let us consider the following cases which are exhaustive by symmetry. If x sees three vertices of A inducing a P_3, then one has a contradiction to Proposition 1. If x sees two vertices of left(A), say v_1, v_2, and a vertex of right(A) missing v_1, v_2, that is vertex v_6, then x, v_1, v_2, v_3, v_4 induce a dart. $\diamond$

Since G is hole-free, we have:

Claim 3.4 *For all $x, y \in A^+$ with x missing y, we have $N_A(x) \subseteq N_A(y)$ or vice versa. Moreover, x and y have a common neighbor in Q_v.*

Claim 3.5 *If $A_2^+ \neq \emptyset$ then $A_3^+ = \emptyset$ and vice versa.*

Proof of Claim 3.5. Assume not; let $x \in A_2^+$ and $y \in A_3^+$, say $N_A(x) = \{v_1, v_4\}$ and $N_A(y) = \{v_1, v_2, v_3\}$ by Claims 3.2 and 3.3. Then since y, v_3, v_6, v_4, x induce no C_5, x misses y but now, by Claim 3.4, the neighborhoods of x and y must be comparable - contradiction. $\diamond$

Claim 3.6 *At most one of $A_2^+(1,4)$, $A_2^+(2,5)$, $A_2^+(3,6)$ is nonempty.*

Proof of Claim 3.6. Assume not; without loss of generality, let $x \in A_2^+(1,4)$ and $y \in A_2^+(2,5)$. Then by Claim 3.4, x sees y but now x, y, v_2, v_3, v_6, v_4 induce a C_6. $\diamond$

Claim 3.7 *The set $N_A(A_1^+)$ of neighbors of all $x \in A_1^+$ in A is a clique.*

Proof of Claim 3.7. Assume not; let without loss of generality $x \in A_1^+$ see v_1 and $y \in A_1^+$ see v_5. Then since x, v_1, v_2, v_5, y do not induce a C_5, x misses y but now by Claim 3.4, $N_A(x)$ and $N_A(y)$ must be comparable - contradiction. $\diamond$

Claim 3.8 *No vertex in A_1^+ sees a vertex in $A_2^+ \cup A_3^+$, i.e., A_1^+ has a co-join to $A_2^+ \cup A_3^+$.*

Proof of Claim 3.8. Assume not; let $x \in A_1^+$ and first assume that $y \in A_2^+$, say $N_A(y) = \{v_1, v_4\}$ with x seeing y. If x sees v_1, then x, v_1, y, v_4, v_2 induce a dart, and if x sees v_2 then x, y, v_2, v_5, v_4 induce a C_5. The other cases are symmetric. This shows Claim 3.8. $\diamond$

Claim 3.9 A_3^+ *is a clique.*

Proof of Claim 3.9. First note that $A_3^+(1, 2, 3)$ has a join to $A_3^+(4, 5, 6)$ since G is C_5-free. If there are $x, y \in A_3^+(1, 2, 3)$ with x missing y then v_4, v_1, v_2, x, y induce a dart. Then $A_3^+(1, 2, 3)$ is a clique. The fact that $A_3^+(4, 5, 6)$ is a clique is shown analogously. $\diamond$

Now we conclude that in any case, we get a clique separator between Q_v and some vertex in A (which finally contradicts to the assumption that G is an atom):

Case 1. $A_3^+ \neq \emptyset$.

Then by Claim 3.5, $A_2^+ = \emptyset$. First suppose that $A_1^+ \neq \emptyset$. We claim that $N_A(A_1^+) \cup A_3^+$ is a clique separator: Recall that by Claim 3.7, $N_A(A_1^+)$ is a clique, by Claim 3.9, A_3^+ is a clique, and by Claims 3.4 and 3.8, every $x \in A_3^+$ sees every $y \in N_A(A_1^+)$ (note that if $N_A(A_1^+) \subseteq \{v_1, v_2, v_3\}$ then $A_3^+(4, 5, 6) = \emptyset$, and the case that $N_A(A_1^+)$ is one of the matching edges is impossible if $A_3^+ \neq \emptyset$). Obviously, $N_A(A_1^+) \cup A_3^+$ is a separator between v and a nonempty part of A.

Now suppose that $A_1^+ = \emptyset$. Then A_3^+ is a clique separator (in this case, possibly $A_3^+(1, 2, 3) \neq \emptyset$ and $A_3^+(4, 5, 6) \neq \emptyset$).

Case 2. $A_2^+ \neq \emptyset$.

Then by Claim 3.5, $A_3^+ = \emptyset$, and by Claim 3.6, at most one of the sets $A_2^+(1, 4)$, $A_2^+(2, 5)$, $A_2^+(3, 6)$ is nonempty, say $A_2^+(1, 4) \neq \emptyset$ and $A_2^+(2, 5) = A_2^+(3, 6) = \emptyset$. Then $\{v_1, v_4\}$ is a clique separator (note that in this case, by Claims 3.4 and 3.8, $N_A(A_1^+) \subseteq \{v_1, v_4\}$).

Case 3. $A_2^+ \cup A_3^+ = \emptyset$.

Then $A_1^+ \neq \emptyset$ since G is connected, and again by Claim 3.7, $N_A(A_1^+)$ is a clique separator. This finishes the proof of Lemma 3. $\square$

By Lemmas 2 and 3, one obtains the following result:

Theorem 7 (*Hole, dart*)*-free atoms are nearly weakly chordal.*

By Theorem 5 (i), we obtain:

Corollary 4 *The MWIS problem is solvable in polynomial time for* (*hole, dart*)*-free graphs.*

The corresponding result in [1] gives a time bound of $\mathcal{O}(n^4)$ for the MWIS problem on (hole, dart)-free graphs. This result is based on Theorem 4 in [1] showing that (hole,dart)-free atoms are nearly gem-free. Their time bound is better than the obvious time bound resulting from Theorem 7 and Theorem 5 (ii). However, more can be said about the structure of (hole,dart)-free graphs: Lemma 2 in [4] implies that prime dart- and gem-free graphs are diamond-free. This is slightly better than Lemma 1 in [1] saying that prime dart- and gem-free graphs are paraglider-free.

Moreover, we need the following result:

Theorem 8 ([3]) (*Hole,diamond*)-*free atoms that contain no induced* $\overline{C_6}$ *are either a clique or chordal bipartite.*

Theorem 8 and Lemma 3 imply:

Corollary 5 (*Hole, dart*)-*free prime atoms are either nearly a clique or nearly chordal bipartite.*

This does not improve the $\mathcal{O}(n^4)$ time bound of [1] but gives more structural insight as asked for in the last paragraph of [1].

4 Structure and MWIS for (Odd-)Hole- and Bull-Free Graphs

In this section, we show that prime odd-hole- and bull-free graphs are nearly perfect and that prime hole- and bull-free atoms are nearly weakly chordal. This is the main result of the section and implies polynomial time for MWIS on both graph classes (with better time bound for (hole, bull)-free graphs). We first collect some properties of prime bull-free graphs.

Let G be a prime graph with at least 7 vertices, and suppose that G contains a $\overline{C_k}$, say H, for some $k \geq 6$. Since G is prime, H is not a module, and thus, there is a vertex $z \notin V(H)$ distinguishing vertices $x, y \in V(H)$, say $xz \in E$ and $yz \notin E$. Since H is connected, we can assume without loss of generality that z distinguishes an edge $xy \in E$ in H. Let $H_0 := H$ and for $k \geq 1$, let H_k result by adding a distinguishing vertex $z_k \notin V(H_{k-1})$ to H_{k-1}, that is, $H_k := G[V(H_{k-1}) \cup \{z_k\}]$. As before, we can assume without loss of generality that z_k distinguishes an edge in H_{k-1}. This defines a strictly increasing sequence of induced subgraphs of G. Since G is finite, there is a largest $\hat{k}$ such that $H_{\hat{k}}$ exists.

Lemma 4 *If G is a connected bull-free graph containing some induced subgraph H_k, $k \geq 0$, as defined above, with nonempty anti-neighborhood $A(H_k)$, and $u \in H_k^+$ then u has a join to H_k.*

Proof. We show Lemma 4 by induction on k. For $k = 0$, the proof goes as follows. Let $u \in H_0^+$ and let $v \in A(H_0)$ be a neighbor of u. To prove the assertion, let us show that if u sees a vertex v_i for some index i, then u sees also v_{i-1} and v_{i+1} (index arithmetic modulo k); by iterating this argument the assertion follows.

Then let us assume without loss of generality that u sees v_1, and show that u sees v_k and v_2. To avoid that $u, v_1, v_3, v_{k-1}, v_k$ induce a bull, u sees either v_3 or v_{k-1} or v_k.

Moreover, by a symmetry argument, one has that u sees either v_3 or v_{k-1} or v_2. Also, if u sees v_k and v_2, then u sees v_{k-1} and v_3, otherwise v, u, v_k, v_2, v_{k-1} or v, u, v_k, v_2, v_3 induce a bull. All the previous facts imply that u sees either v_{k-1} or v_3; without loss of generality, say u sees v_3. Then u sees v_k, otherwise v, u, v_1, v_3, v_k induce a bull. Then u sees v_2, otherwise v, u, v_3, v_k, v_2 induce a bull. This completes the proof for $k = 0$.

Now, as the induction hypothesis, assume that the assertion is true for values at most k, and let us show that it holds for $k + 1$. Let $z \notin V(H_k)$ be any vertex distinguishing an edge xy in H_k, say $xz \in E$ and $yz \notin E$, and let $H_{k+1} = G[V(H_k) \cup \{z\}]$. Suppose that $A(H_{k+1}) \neq \emptyset$, and without loss of generality choose a vertex $v \in A(H_{k+1})$ such that the distance between v and H_{k+1} is 2. Now, also $v \in A(H_k)$ holds, and the distance between v and H_k is at least 2. If the distance is 2 then let $u \in H_k^+$ be a neighbor of v. Thus, v misses H_k and, by the induction hypothesis, any neighbor $u \in H_k^+$ of v has a join to H_k. Since $v \in A(H_{k+1})$, $vz \notin E$ holds. Since G is bull-free, v, u, x, y, z is not a bull, and since v misses x, y, z and u sees x, y, it follows that $uz \in E$ and thus, u has a join to H_{k+1}.

In the other case, assume that the distance between v and H_k is at least 3, and recall that the distance between v and H_{k+1} is 2; thus, let u be a common neighbor of v and z. Note that in this case, u misses H_k, and now by the induction hypothesis, $z \in H_k^+$ must have a join to H_k which is a contradiction. This shows Lemma 4. □

Lemma 5 *Prime bull-free graphs are nearly $\overline{C_\ell}$-free for any $\ell \geq 6$.*

Proof. Suppose that there is a vertex v such that $A(v)$ contains a $\overline{C_\ell}$, say H, for some $\ell \geq 6$. Without loss of generality, let v be in distance 2 to H, and let $u \in H^+$. Thus, by definition and by Lemma 4, H is contained in $A(v) \cap N(u)$. Now let H_k, $k \geq 0$, and $H_{\hat{k}}$ be defined as above. We claim that v misses $H_{\hat{k}}$: Clearly v misses H_0. Suppose that v misses H_k and sees H_{k+1}; recall that H_{k+1} results by adding a distinguishing vertex z_k to H_k; say, for some edge xy in H_k, z_k sees x and misses y. Now if v misses H_k and sees H_{k+1} then $vz_k \in E$, and now, since $z_k \in H_k^+$, by Lemma 4, z_k should be universal for H_k - a contradiction.

Thus, by Lemma 4, it follows that $H_{\hat{k}}$ is a module in G which is a contradiction to the assumption that G is prime. □

By the Strong Perfect Graph Theorem and by the definition of weakly chordal graphs, it follows:

Corollary 6 *Prime (odd-hole,bull)-free graphs are nearly perfect, and prime (hole,bull)-free graphs are nearly weakly chordal.*

By Theorem 2, Corollary 1 and Theorem 5, we obtain:

Corollary 7 *For (odd-hole,bull)-free graphs, the MWIS problem is solvable in polynomial time.*

The time bound for (hole,bull)-free graphs is $\mathcal{O}(n^5)$.

Note that Corollary 7 concerning (odd-hole,bull)-free graphs is close to the MWIS result implied by the structure result of De Simone [16], i.e., MWIS is polynomial for graphs with no odd apples and no six specific induced subgraphs, five of which contain

a bull. It is also well known that MWIS can be solved in polynomial time for P_5- and bull-free graphs since such graphs do not contain any of the forbidden subgraphs in [16]. However, for P_5- and bull-free graphs, we can say more; the proof of Lemma 4 can be adapted to similar cases as we describe subsequently.

5 MWIS for P_5- and Bull-Free Graphs Revisited

Lemma 6 *Prime P_5- and bull-free graphs are nearly C_5-free and nearly house-free.*

Proof. We recursively define a sequence H_k, $k \geq 0$, of increasing subgraphs for which $A(H_k)$ is nonempty, $v \in A(H_k)$ and $u \in H_k^+$ being a neighbor of v in the same way as before Lemma 4; only H_0 is different. For showing that prime P_5- and bull-free graphs are nearly C_5-free, H_0 is a C_5, say with vertices $v_1, \ldots, v_5$ and edges $v_i v_{i+1}$ (index arithmetic modulo 5), and for showing that prime P_5- and bull-free graphs are nearly house-free, H_0 is a house, say with vertices $v_1, \ldots, v_5$ such that v_1, v_2, v_3, v_4 induce a C_4 and v_5 sees v_2 and v_3.

We show Lemma 6 by induction on k. Let G be a prime (P_5,bull)-free graph. For $k = 0$, we first consider the case when H_0 is a C_5. We need the following notion:

For a subgraph H of G, a vertex $x \notin V(H)$ is an *i-vertex of H* if x sees exactly i vertices in H. Obviously, C_5 has no 1-vertex since G is P_5-free, and C_5 has no 2-vertex since two consecutive neighbors lead to a bull, and two nonconsecutive neighbors lead to P_5. Similarly, if for a 3-vertex, not all neighbors are consecutive, we get a bull, and for consecutive neighbors, we get a P_5. Any 4-vertex leads to a bull. Thus, any vertex outside $H_0 = C_5$ seeing H_0 is universal for H_0. Now with the same inductive arguments as in the proof of Lemma 4, we show that any vertex $u \in H_k^+$ has a join to H_k. With the same arguments as in Lemma 5, we obtain that G is nearly C_5-free.

Now, let H_0 be a house as described above.

1. For a 1-vertex u, we obtain a P_5 if u is adjacent to v_1 or v_4 or a bull if u is adjacent to v_2, v_3 or v_5.

2. For a 2-vertex u, there are two cases:

    (a) u has two neighbors in the C_4:
        i. if u sees v_1 and v_2 then u, v_1, v_2, v_4, v_5 induce a bull,
        ii. if u sees v_1 and v_3 then v, u, v_1, v_2, v_5 induce a P_5,
        iii. if u sees v_2 and v_3 then v, u, v_2, v_3, v_4 induce a bull, and
        iv. if u sees v_1 and v_4 then v, u, v_1, v_3, v_4 induce a bull.
    (b) u has one neighbor in the C_4 and sees v_5:
        i. if u sees v_1 then u, v_2, v_3, v_4, v_5 induce a bull, and
        ii. if u sees v_2 then v, u, v_1, v_2, v_5 induce a bull.

3. For a 3-vertex u, there are two cases:

    (a) u has three neighbors in the C_4:

     i. if u sees v_1, v_2 and v_3 then u, v_1, v_2, v_4, v_5 induce a bull, and

     ii. if u sees v_1, v_2 and v_4 then v, u, v_1, v_3, v_4 induce a bull.

(b) u has two neighbors in the C_4 and sees v_5:

     i. if u sees v_1 and v_2 then v, u, v_1, v_2, v_4 induce a bull,

     ii. if u sees v_1 and v_3 then v, u, v_3, v_4, v_5 induce a bull,

     iii. if u sees v_2 and v_3 then v, u, v_2, v_3, v_4 induce a bull, and

     iv. if u sees v_1 and v_4 then v, u, v_1, v_3, v_4 induce a bull.

4. For a 4-vertex u, there are two cases:

(a) If u has four neighbors in the C_4 then v, u, v_3, v_4, v_5 induce a bull.

(b) Otherwise, if u has three neighbors in the C_4 and sees v_5 then:

     i. if u sees v_1, v_2 and v_3 then v, u, v_1, v_2, v_4 induce a bull, and

     ii. if u sees v_1, v_2 and v_4 then v, u, v_1, v_3, v_4 induce a bull.

Thus, u is universal for H_0, and with the same arguments as above, we obtain that G is nearly house-free.   □

A bipartite graph B with color classes X and Y is a *bipartite chain graph* if the neighborhoods of the vertices of one color class form an increasing sequence with respect to set inclusion. It is well known that these graphs can be recognized in linear time, have bounded clique-width, and MWIS can be solved in linear time for them. From a result by Fouquet [21], it follows:

Corollary 8 *Prime (P_5,bull)-free graphs are nearly bipartite chain graphs or nearly co-bipartite chain graphs.*

Corollary 9 *For (P_5,bull)-free graphs, the MWIS problem is solvable in time $\mathcal{O}(nm)$.*

A polynomial time algorithm for MWIS on (P_5,dart)-free graphs follows from [25].

6   Conclusion

The class of hole-free graphs is closely related to many well-studied graphs classes, such as chordal graphs, weakly chordal graphs and perfect graphs. However the complexity of the MWIS problem for hole-free graphs is open. Following a recent line of research [5, 6], we show that MWIS can be solved in polynomial time for odd-hole- and dart-free graphs as well as for odd-hole- and bull-free graphs. For this purpose, the main structural results of this paper are the following (which might also be useful in other contexts):

1. (a) Dart-free graphs are nearly $\overline{C_{2k+1}}$-free for $k \geq 3$.

   (b) Hole- and dart-free atoms are nearly $\overline{C_k}$-free for $k \geq 6$.

   (c) Consequently, odd-hole- and dart-free graphs are nearly perfect, and hole- and dart-free atoms are nearly weakly chordal.

2. (a) Prime bull-free graphs are nearly $\overline{C_k}$-free for any $k \geq 6$.

   (b) Consequently, prime odd-hole- and bull-free graphs are nearly perfect, and prime hole- and bull-free atoms are nearly weakly chordal.

3. Using the approach for bull-free graphs, we show that prime (P_5,bull)-free graphs are nearly bipartite chain graphs or nearly co-bipartite chain graphs; this leads to a better time bound for MWIS on (P_5,bull)-free graphs.

The results on dart-free graphs imply that MWIS is solvable in polynomial time for odd-hole- and dart-free graphs, by finally reducing the problem to perfect graphs. Note that (hole, dart)-free graphs can be recognized in polynomial time since recognition of hole-free graphs can be done in polynomial time by [26] (see also [13] for recognizing dart-free perfect graphs in polynomial time, though in general perfect graphs can be recognized in polynomial time by [11, 15]).

The results on bull-free graphs allow to solve MWIS in polynomial time for odd-hole- and bull-free graphs, again by finally reducing the problem to perfect graphs. Note that (hole, bull)-free graphs can be recognized in polynomial time since recognition of hole-free graphs can be done in polynomial time by [26].

Let us conclude by a remark on the class of (hole,gem)-free graphs. It seems that the situation for (hole,gem)-free graphs is more complicated than for (hole,dart)-free and (hole,bull)-free graphs. Clearly, by the Strong Perfect Graph Theorem, (hole,gem)-free graphs are perfect, and thus, the MWIS problem is solvable in polynomial time for (hole,gem)-free graphs. However, we would like to find a direct combinatorial algorithm with a good time bound as in the other cases. Therefore we conclude with the following:

Open Problem. What is a good time bound for the MWIS problem for (hole,gem)-free graphs?

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